



DESIGNING DETECTION OF DATA LINEAGE FRAMEWORK IN ADVERSARIAL ENVIRONMENTS

PATAN NAYEEMULLAH KHAN PG Scholar, Dept of CSE, Quba College of Engineering & Technology, Venkatachalam, AP, India.

D.MURALI Assistant Professor, Dept of CSE, Quba College of Engineering & Technology, Venkatachalam, AP, India.

ABSTRACT

A growing number of customers are considering cloud computing as prospective hosts for their enormous datasets. Existing research recommends encrypting sensitive data before outsourcing and implementing Searchable Symmetric Encryption (SSE) to enable keyword-based searches over the ciphertexts because the cloud service provider (CSP) is outside the users' trusted domain. Designing a successful SSE system that concurrently supports sublinear search time, effective update and verification, and on-demand information retrieval is still a difficult issue.. To solve this, we provide the Verifiable Dynamic Encryption with Ranked Search (VDERS) method, which enables a user to securely and quickly perform top-K searches on a dynamic document collection and confirm the accuracy of the search results. To be more precise, we first offer a fundamental structure, VDERS0, that enables the insertion of verified documents in top-K searches by creating a ranking inverted index and a verifiable matrix. Subsequently, VDERS, a sophisticated architecture, is developed to facilitate document deletion even further while requiring less communication. Many tests on real datasets show how efficient and effective our VDERS system is.

1. INTRODUCTION

Cloud computing has received a lot of attention from the research and business industries as a potential computing paradigm. Users frequently outsource their big datasets to clouds and assign a cloud service provider (CSP) to take care of data storage and provide query services because of the advantages of low prices, flexibility, and scalability. Research currently available advises encrypting data before outsourcing due to security and privacy issues [1]. Yet, keyword-based searches over ciphertexts are a difficult problem due to data encryption. In a dynamic and malevolent cloud environment, this is considerably more difficult [2].

Let's think about the next circumstance. In the cloud, where each email is indexed by the sender's name and sorted by receipt date, Alice outsources her archived emails. For instance, the email received on April 2 has a higher rank than the email received on April 1 for a set of emails indexed by the term Bob. Alice uploads her documents and keywords to the cloud in encrypted form to protect their confidentiality. There could be hundreds of documents that match a given keyword, and the user would incur significant fees if they were all returned to them and decrypted. Alice may therefore wish to run a top-K search to find the most current emails..



Also, Alice might like to save money by just storing emails that she has received in the recent three months. For instance, Alice will delete all emails she received prior to February whenever May arrives. The chosen encryption method should adhere to the following specifications in the aforementioned application scenario: First, a ranked search. To find the papers that are most closely matched, the user is permitted to conduct a top-K search. (2) Active. Documents kept in the cloud can be updated (added and removed) by the user. (3) Validity. To conserve memory, the malicious CSP may destroy encrypted files that are not frequently accessed or it may falsify search results to trick the user..

A virus or worm can tamper with encrypted documents even if the CSP is truthful. As a result, the user should be able to confirm the accuracy of the search results. 4. Effectiveness. A set of encrypted documents may be efficiently searched, updated, and verified by the user..

Existing SSE schemes only partially meet the aforementioned requirements, despite the fact that Searchable Symmetric Encryption (SSE) enables users to access requested documents while protecting their anonymity. This work suggests a Verifiable Dynamic Encryption with Ranked Search (VDERS) method that enables the user to carry out updates and top-K searches on ciphertexts in a verifiable and effective manner in order to simultaneously satisfy all these features. The key concept behind our approach is to build a verifiable matrix to store the rating data and encrypt it using an RSA accumulator [3]. Furthermore, to enable effective top-K searches and updates, a ranking inverted index is created from a set of documents..To be more precise, we first offer a fundamental architecture, indicated by VDERS0, that allows for verified document insertion activities. Then, we offer a sophisticated design, abbreviated as VDERS, that not only enables effective deletion operations but also lowers communication costs without outsourcing the verifiable matrix. These are our primary contributions, in brief:

To enable dynamic and ranked searches in a cloud environment in an effective and verifiable manner, we suggest a VDERS scheme.To achieve effective top-K searches with support for verifiable updates, two architectures are offered.To demonstrate the success of our system, we carry out comprehensive experiments on real datasets and theoretically examine its performance and security.

2. PROBLEM STATEMENT

Since Song et al. initially suggested SSE, it has been the subject of extensive research. as a key contribution to SSE. A thorough security definition and two systems, SSE-1 and SSE-2, based on an inverted index, were developed by Curtmola et al. A user can quickly and securely update the encrypted outsourced data using Dynamic SSE (DSSE). With Verifiable SSE (VSSE), a user can confirm the accuracy of search results. A user can spot any dishonest server's cheating behaviour using the Universally Composable (UC)-secure VSSE technique developed by Kurosawa and colleagues. According to the quantity of matching query keywords, Cao et al dynamic 's multi-keyword RSSE system ranked results using the secure KNN technique..

LIMITATION OF SYSTEM



Their plan's primary drawback is its inefficiency. The amount of documents increases linearly with the search time..

3. PROPOSED SYSTEM

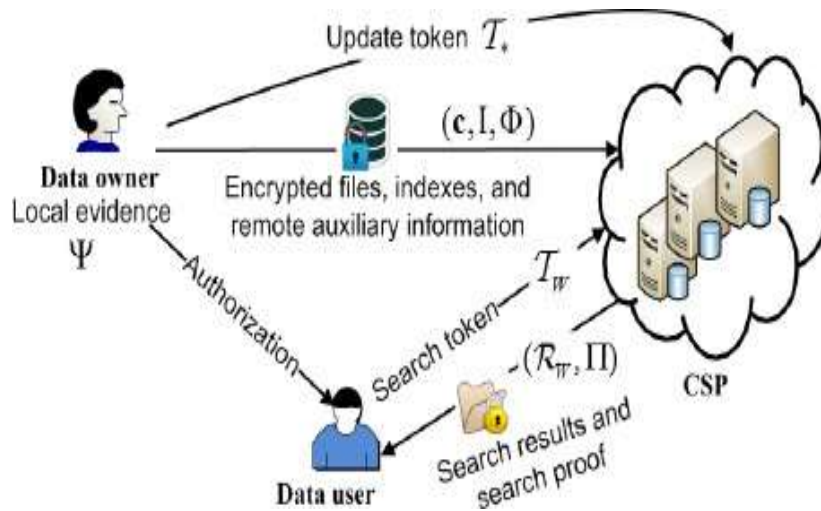
In order to achieve dynamic and ranked searches in a cloud environment in an effective and verifiable manner, we suggest a VDERS scheme. To obtain effective top-K searches with support for verifiable updates, two architectures are offered. To demonstrate the success of our system, we carry out comprehensive experiments on real datasets while theoretically analysing its performance and security..

ADVANTAGES OF PROPOSED SYSTEM

Their plan's primary drawback is its inefficiency. The amount of documents increases linearly with the length of the search. Designing SSE methods that facilitate dynamic update for cloud data is essential.

4. SYSTEM IMPLEMENTATION

As shown in Fig. 1, the system is composed of three distinct parties: the CSP, the data owner, and the data consumer. Data storage and query services are provided by the CSP using cloud platforms that pool hard and soft resources. Cryptotexts are first created by the data owner for a document collection d . She then constructs a secure index I for quick searches using the keywords w derived from d , and produces local evidence and distant auxiliary information for verifiable queries. She can update ciphertexts with an update token T after uploading $(c, I,)$ to the cloud and retrieve documents whenever she needs them with a search token T in a verifiable manner. After confirming the accuracy of the search results, the data owner recovers document contents after getting the search results and a search proof $(RW,)$ from the CSP. The data owner can also assign authorised data users the responsibility for searching, updating, and verifying the data. In this work, we refer to both the data owner and the data user together as "users".



5. IMPLEMENTATIONALGORITHM

Algorithm 1 VDERS⁰.Encryption

Input: Secret key SK , documents \mathbf{d} , and keywords \mathbf{w}

Output: Ranked inverted index I and ciphertexts \mathbf{c}

- 1: Initialize the search table T_s with an empty dictionary of size $m + 1$ and initialize the search array A_s with an empty array of size $\|\mathbf{c}\|/8 + z$
 - 2: **for** each keyword $W \in \mathbf{w}$ **do**
 - 3: Create a ranked linked list $L_W = (N_1, \dots, N_{|\mathbf{d}_W|})$ by randomly choosing $|\mathbf{d}_W|$ unused entries in A_s where $N_i = \langle id_i, \varepsilon_i, addr_s(N_{i+1}) \rangle$ is defined in Definition 1
 - 4: **for** $i \in [|\mathbf{d}_W|]$ **do**
 - 5: Choose a κ -bit random string r_i and set $A_s[addr_s(N_i)] \leftarrow (N_i \oplus \mathbb{H}(P_{k_3}(W), r_i), r_i)$
 - 6: Encrypt the address of the head node N_1 and set $T_s[F_{k_1}(W)] \leftarrow addr_s(N_1) \oplus G_{k_2}(W)$
 - 7: Create a free list $L_{free} = (F_1, \dots, F_z)$ by randomly choosing z unused entries in A_s where $F_i = \langle \mathbf{0}, \mathbf{0}, addr_s(F_{i-1}) \rangle$ and $addr_s(F_0) = -1$
 - 8: **for** $i = z$ to 1 **do**
 - 9: Set $A_s[addr_s(F_i)] \leftarrow (F_i, \mathbf{0})$
 - 10: Store the address of the tail node F_z in T_s by setting $T_s[free] \leftarrow addr_s(F_z)$
 - 11: Fill the remaining entries of A_s with random strings
 - 12: Set $I = (T_s, A_s)$
 - 13: **for** $j \in [n]$ **do**
 - 14: Set $C_j = \text{SKE.Enc}(k_e, D_j)$ and $\mathbf{c} = \mathbf{c} \cup C_j$
 - 15: **return** (I, \mathbf{c})
-

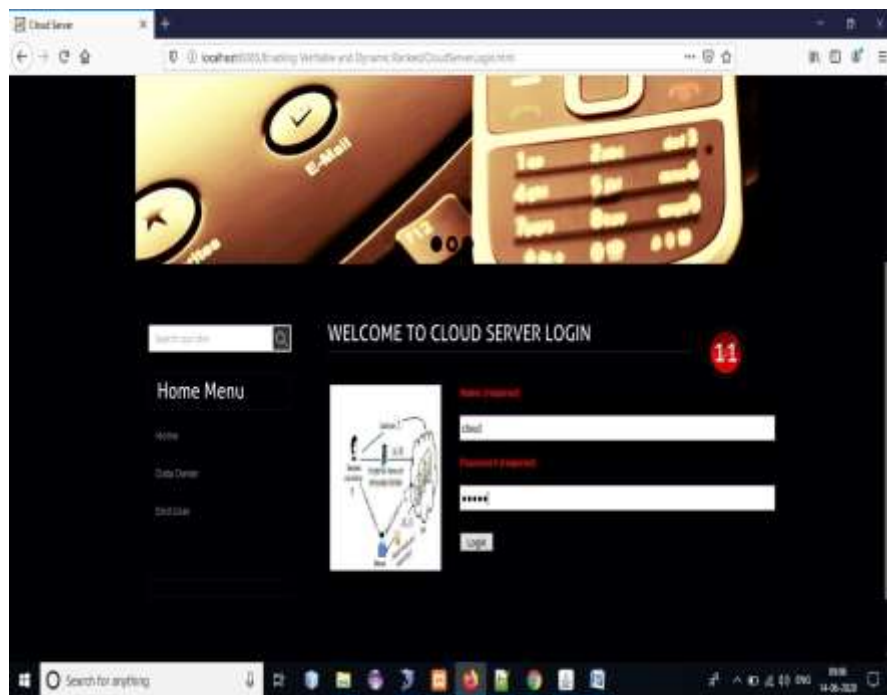
Algorithm 2 VDERS⁰.Search

Input: Ranked inverted index I and search token \mathcal{T}_W

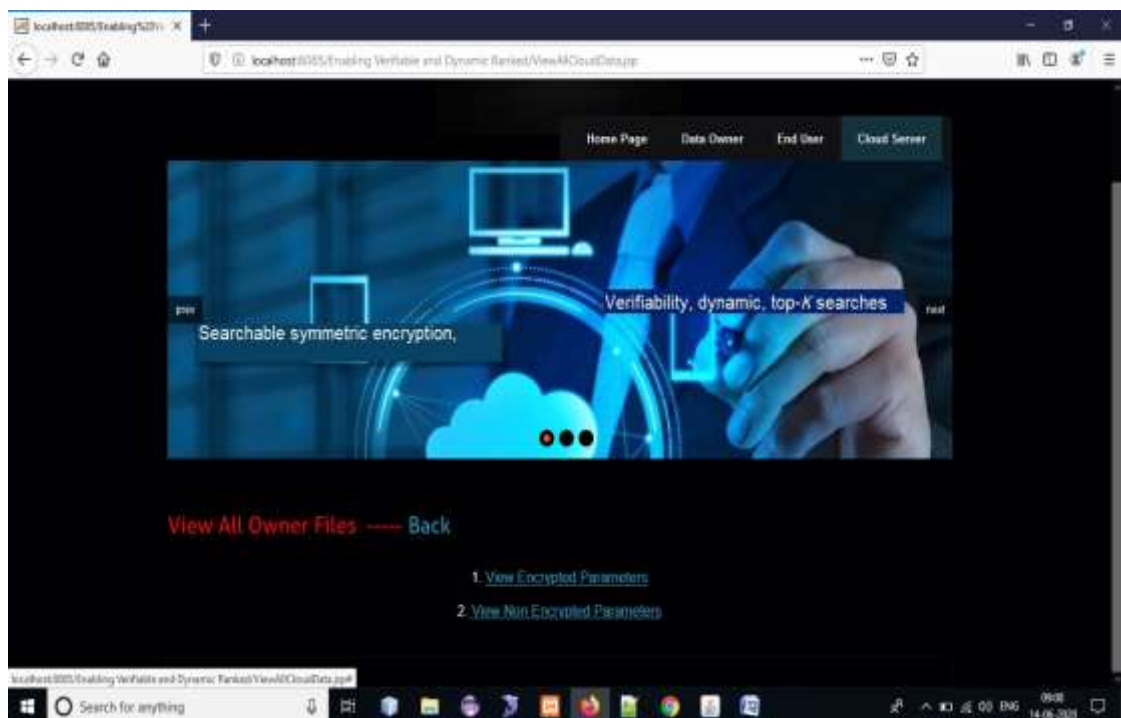
Output: Search results \mathcal{R}_W

- 1: Parse \mathcal{T}_W as $(\tau_1, \tau_2, \tau_3, \tau_4, \tau_5)$
 - 2: **if** τ_1 is not in T_s **then**
 - 3: **return** \emptyset
 - 4: Compute $T_s[\tau_1] \oplus \tau_2$ and recover $addr_s(N_1)$ the address of the head node N_1 of L_W in A_s
 - 5: **repeat**
 - 6: Obtain the encrypted node N_i stored at $A_s[addr_s(N_i)]$ and parse N_i as (v_i, r_i)
 - 7: Compute $v_i \oplus H(\tau_3, r_i)$ and recover node $N_i = \langle id_i, \varepsilon_i, addr_s(N_{i+1}) \rangle$
 - 8: **until** $addr_s(N_{i+1}) = -1$
 - 9: Sort documents according to their scores
 - 10: Let ID_j and ε_j be the identifier and score of the rank- j document for W , respectively
 - 11: **if** $\tau_4 \neq *$ **then**
 - 12: Set $id_W^{K+1} = (ID_1, \dots, ID_{K+1})$, $s_W^{K+1} = (\varepsilon_1, \dots, \varepsilon_{K+1})$, and $c_W^K = (C_{ID_1}, \dots, C_{ID_K})$
 - 13: Set $\mathcal{R}_W \leftarrow (id_W^{K+1}, s_W^{K+1}, c_W^K)$
 - 14: **else**
 - 15: Set $id_W = (ID_1, \dots, ID_{|d_W|})$, $s_W = (\varepsilon_1, \dots, \varepsilon_{|d_W|})$, and $c_W = (C_{ID_1}, \dots, C_{ID_{|d_W|}})$
 - 16: Set $\mathcal{R}_W \leftarrow (id_W, s_W, c_W)$
 - 17: **return** \mathcal{R}_W
-

6. OUTPUTSCREENS







View Attribute Based Encrypted Parameters ----- Back

Index Id	File Name	Trapdoor	Secret Key	OwnerName	Search Token
7	Connect.jsp	-4219fab95dc0d-902d629046662baf936193b41	JB@127010f	UmfqZXNo	Over Connection
8	DotNet.txt	LTNed4M3MTNwYjMwZDZlYmY2NmY3MmY0YTs2DQZTU4NjNYQ==	W5A2Thy0TE==	UmfqZXNo	Dot Net Programms
9	Java.txt	LTFYtZHYT2m4chNkxM1crYzJkNzE4ZDZmRGlUMQlKS2TA1M2l1YmQ=	W5AMT0BMJ-52A==	UmfqZXNo	Java Programms
10	EBank.txt	MTBjCOMNw1MwA7NzCC052mQ1YjZ3NTYTMWQ2SM0YxODY5Nm==	WQJAMWJZDwMg==	UmfqZXNo	Banking Application
11	Ecommerce.txt	Mdc2mMYMCRyYADYTBQzQzDMZYTtR0DQ1Y2NTYzJkwZjg=	WUJAMW1lyYmMmYQ==	UmfqZXNo	Ecommerce Applications
12	GST.txt	NjNkZmZlZDZmRGlUMQlKS2TA1M2l1YmQ=	54325343	TWFuamV4YXRo	GST Filing Techniques
13	Java.txt	LTFYtZHYT2m4chNkxM1crYzJkNzE4ZDZmRGlUMQlKS2TA1M2l1YmQ=	WUJAMT13Z3E-1ZQ==	TWFuamV4YXRo	OOPS
14	Android.txt	-8-1806a20203d6680d96837510301beek2ue	JB@30670	TWFuamV4YXRo	Android App
16	auth.jsp	WTFNzU1NDk2YTZmQzQzZmZlZDZmRGlUMQlKS2TA1M2l1YmQ=	WUJAMW1Z3TRMw==	TWFuamV4YXRo	Cloud Authentication
17	Bigdata.txt	T2a78Ba1d8c1682a4Hb9H79537443838008	JB@R33a0y	52WYXh	Big Data Processing

WELCOME TO DATA OWNER REGISTRATION

11

Search results

Sidebar Menu

- Home
- End User
- Cloud Server

User Name (required)

zree

Password (required)

Email address (required)

zree@gmail.com

Mobile Number (required)

9903843003

Your Address

9k11.com

Date of Birth (required)

11/02/1995

Sex (Gender) (required)

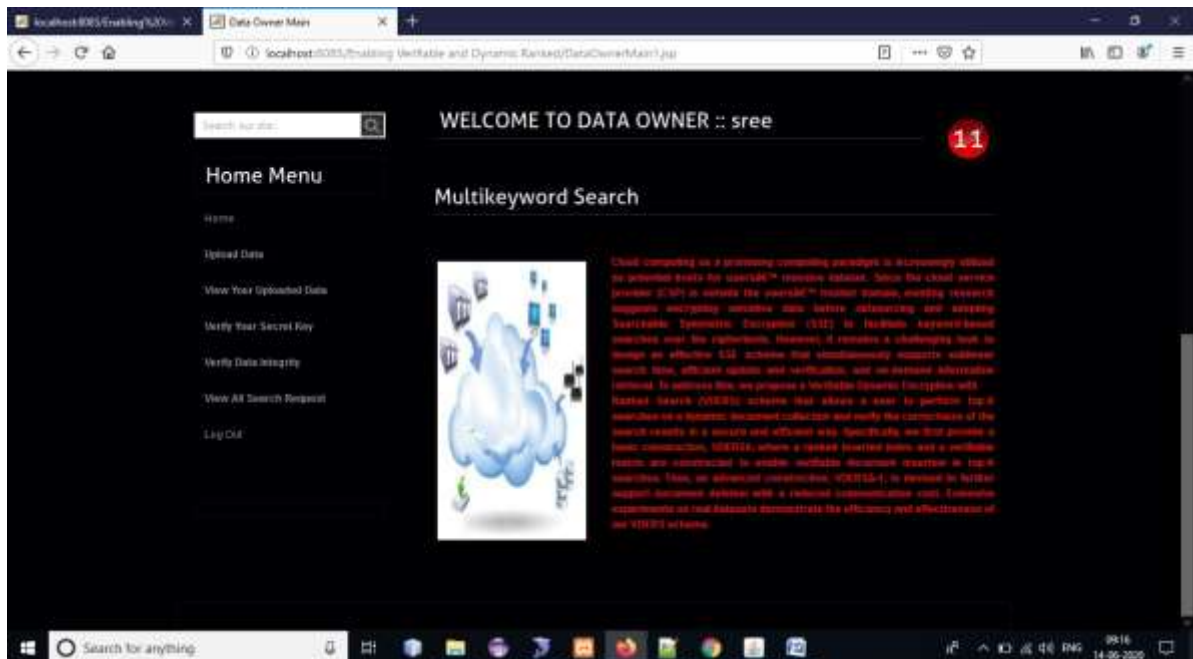
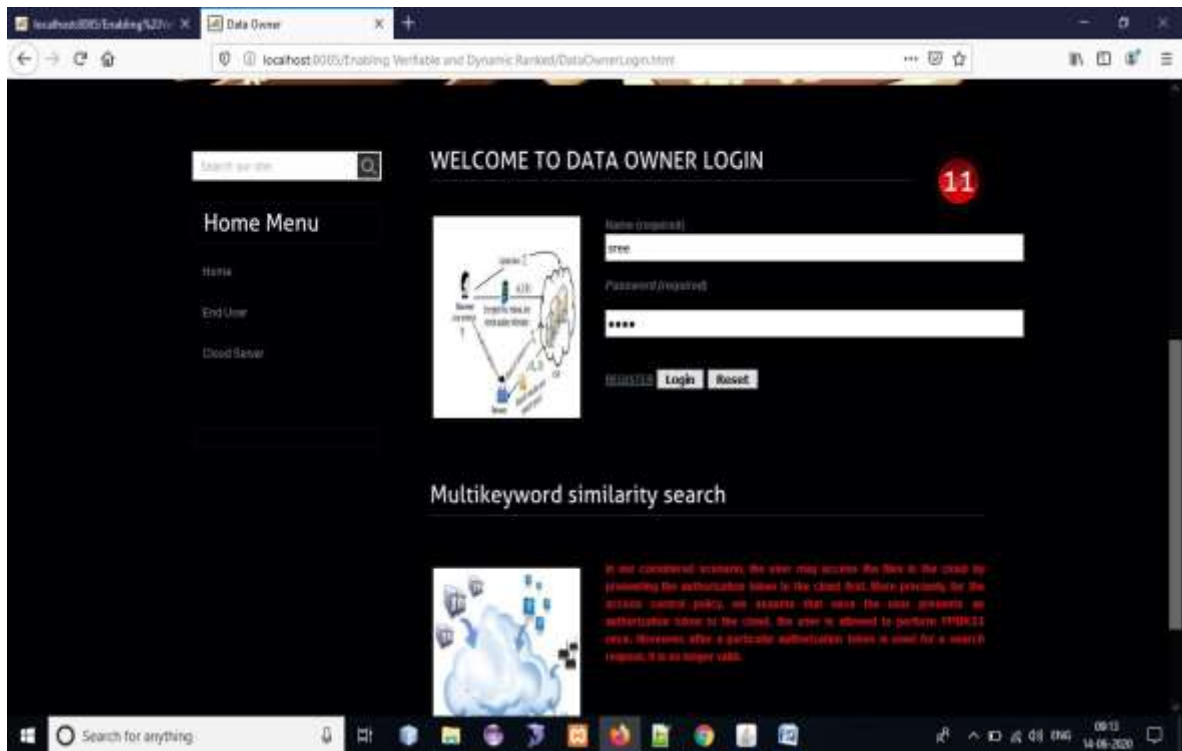
MALE

Enter Pincode (required)

524001

Enter Location (required)

Hydore



7. CONCLUSIONS AND DISCUSSIONS

In this research, we build a VDERS technique to simultaneously allow on-demand information retrieval in a cloud computing environment, efficient updating and verification, and sublinear search time. Results from experiments show how useful our method is for confirming the accuracy of top-K searches on a dynamic document collection. We will endeavour to create a forward secure VDERS



scheme in the next work that prevents keyword information about newly uploaded documents from being leaked during the update phase..

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