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SIMILARITY SEARCH FOR ENCRYPTED IMAGES IN SECURE CLOUD COMPUTING

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## ABSTRACT:

With the emergence of intelligent terminals, the Content-Based Image Retrieval (CBIR) technique has attracted much attention from many areas (i.e., cloud computing, social networking services, etc.). Although existing privacy-preserving CBIR scheme scan guarantee image privacy while supporting image retrieval, these schemes still have inherent defects (i.e., low search accuracy, low search efficiency, key leakage, etc.). To address these challenging issues, in this paper we provide a similarity Search for Encrypted Images in secure cloud computing (called SEI). First, the feature descriptors extracted by the Convolutional Neural Network (CNN)model are used to improve search accuracy. Next, an encrypted hierarchical index tree by using K-means clustering based on Affinity Propagation (AP) clustering is devised, which can improve search efficiency. Then, a limited key-leakage k-Nearest Neighbor (KNN) algorithm is proposed to protect key from being completely leaked to untrusted image users. Finally, SEI is extended to further prevent image users' search information from being exposed to the cloud server. Our formal security analysis proves that SEI can protectimage privacy as well as key privacy. Our empirical experiments using a real-world dataset illustrate the higher search accuracy and efficiency of SEI.

Keywords: content, based, image, retrieval, CNN, security analysis

## INTRODUCTION :

WITH the rapid development and popularization ofcloud computing, people enjoy various conveniencesbrought by cloud

services, such as storing images on thecloud. However, directly outsourcing images to the

publiccloud inevitably raises privacy concerns. Once the massiveimages (e.g., patients' medical images) containing highlysensitive information have been leaked to unauthorizedensities, it will incur serious consequences or unnecessarytrouble. The encryption mechanism can alleviate image datasecurity and privacy concerns to some extent, but it invalidates the Content-Based Image Retrieval (CBIR) techniqueover ciphertext, and even causes other concerns discussed in the following example.

Example. Alice outsources the encrypted images Coif the local image databaseMto the cloud server. The authenticated Bob generates an encrypted search requestTaccording to query imagemqby using the searchable keyskshared by Alice when he queries images similar tomq. Then,the cloud server searchesCafter receivingTand returnsrelevant search resultsRto Bob. Finally, Bob decrypts theRwith image encryption keykiefrom Alicen this encrypted image search process, the cloud serveris regarded as semi-trusted. Alice stores her images onthis cloud server, and entrusts the cloud server to performsimilarity search tasks. When Bob comes to query Alice'simages, the cloud server will honestly provide Bob withsearch service under Alice's arrangement. The performanceof search services, such as search accuracy and efficiency,will profoundly affect Bob's search experience. Assumingthat Bob is a doctor who relies on search results to diagnose certain patient's condition, the incorrect search results will lead to a wrong diagnose, which endangers the patient'shealth and even life. Besides, the time-consuming searchprocess

will prolong the waiting time of the Bob. If Bob isa mobile image user who requires high real-time responses,such a long search time is hard to bear and easily makessearch time lose timeliness. Moreover, the above searchmechanism still requires Alice to shareskandkiewithBob, which cannot completely protect image privacy. Thisis because we cannot promise that Bob is fully trusted anddoes not share kandkiewith other unauthorized imageusers due to interest incentives in practice.

Fortunately, various schemes related to privacy-preserving CBIR have been studied, like. However, n practice, these schemes still face many challenges (i.e.,low search accuracy, low search efficiency, key leakage, etc.). Specifically, schemesdirectly distributed keysto users, leading to the risk of image users leaking keys, schemes sacrificed accuracy to improve efficiency,and schemes brought a lot of overhead to achieve high security. These challenges limiting theirpractical applications are shown.

As for the first challenge indicated by thereare two primary reasons affecting the search accuracy. One isthe feature descriptors type; the other is the similarity calculation method. As for the former, image feature descriptors are mainly classified into global features (e.g., Local BinaryPattern (LBP), Histogram of Oriented Gradients (HOG), etc.) and local<br>features (e.g., Scale Invariant Feature Scale Invariant Feature Transform (SIFT), Speeded-Up Robust Features (SURF), etc.). The local features can achieve higher search accuracy than global features due to the better robustness. There are also works [9],[10] that combine global features and local features with certain weights to form new features or apply Convolutional Neural Network (CNN) model mimicking humanvisual cognition to extract feature vectors, which achieveacceptable accuracy. For the later, the similarity betweenimages is measured by Euclidean distance, Cosine distance,Hamming distance, and Jaccard similarity coefficient. Especially, Asymmetric Scalar-product-Preserving Encryption (ASPE) algorithm using random numbers and matrices to encrypt feature

vectors can calculate the Euclideandistance of high-dimension space more accurately. At thesame time, other works,using Secure MultipartyComputation (SMC), Homomorphic Encryption (HE) to calculate Euclidean distance can also improve search accuracy,but result in inefficiency due to complex garbled circuitoperations.

In order to support the large-scale image search in actualapplications, it is indispensable to improve thesearch efficiency. As far as we know, a feasible searchscheme can be achieved by constructing the efficient linearindex, inverted index, hash table or tree index. In particular,the hash table and tree index perform better than the other two regarding search efficiency. For example, the two-levelhash table constructed based on the Locality SensitiveHashing (LSH) algorithm greatly speeds up the searchprocess by reducing the dimension of the highdimensionfeature vectors. For tree index, scheme [10] classifies imagesby usingKmeans clustering algorithm and then buildsthe index tree based on clustering results, which narrowsthe search scope and avoids traversing the entire image database during the search phase.

Although some prevalent schemes based on ASPE and index tree achieve high search efficiency without deter-rating search accuracy, they do not solve the keyleakage challenge, the image owner should sharekandkiewith all authorized image users. However, it is impractical and overloaded to establish a secure channel to transmit these keys. To avoid key transmission, the latest scheme focusing on improving the ASPE algorithm can guarantee that's kused to encrypt vectors will not be completely disclosed to any image user. Unfortunately, it does not support image search. Although some image search solutions can avoid leakingk,these solutions still have low search accuracy and search efficiency due to the use of homomorphic encryption and secure multiparty computation

As far as we know, no existing work is dedicated to solving the above three challenges simultaneously. Hence, in this

paper we propose a similarity Search for Encrypted Images in secure cloud computing (SEI) to solve the above challenges. Specifically, we employ the CNN model to extractfeature vectors to improve search accuracy, and then build a hierarchical index tree in a bottom-up manner based on the clustering algorithm to improve search efficiency. Besides,we design an optimized ASPE algorithm, which does notrequire the image owner to shareskwith image users, toachieve limited key-leakage for untrusted image users. Tosummarize, our contributions are set out as follows.

High search accuracy. SEI achieves a high search ac-curacy by using the pre-trained CNN model to extractfeature vectors. The CNN model simulating the human visual perception process can more accurately represent the image content, which makes the similarity measurement between images more accurate and search results more accurate

High search efficiency. SEI uses the optimized K-means clustering algorithm to classify images and con-structs a hierarchical index tree in a bottom-up manner based on the clustering results. Therefore, SEI avoids traversing the entire image database when performing search operation and reduces the search time to sublinear time. Moreover, the Principal Component Analysis(PCA) algorithm is utilized to lower the dimensions of the image feature vectors extracted by CNN model,which further improves search efficiency

Limited key leakage. SEI provides a secure trapdoor generation process with limited keyleakage1, which not only prevents untrusted image users from completely leaking keys privacy but also avoids the online help of the image owner when the image user generates trapdoors locally. Moreover, even if any unauthenticated image user illegally obtains partial information about searchable keys from a certain untrusted user, he cannot generate a valid query trapdoor or decrypt theciphertext of the image owner.

## SYSTEM ARCHITECTURE:



#### METHODOLOGY : Local Feature based Scheme:

For local feature-based scheme, a set of local features are extracted from local regions of an image. One popular approach of content-based image retrieval using local features is called bag-of-words model. In this model, local features are extracted from all images in the database and jointly clustered. The cluster centers are used as cluster identifiers, which form the vocabulary. After clustering, only the identifier of the most similar cluster center is kept for each local feature, which signifies a word. Then the cluster occurrence histogram is created for each image, which can be represented as a vector of occurrence counts of the local features. Through this way, an image can be expressed as a bag of words in a visual vocabulary. By comparing the similarity of histograms, one can retrieve similar images. As mentioned before, EMD distance is used during the search process. Denote the signature of each image in the database  $\mathrm{as} S_t =$  $\left\{ (s_1^{(t)}, w_1^{(t)}), (s_2^{(t)}, w_2^{(t)}), \ldots, (s_m^{(t)}, w_m^{(t)}) \right\}$  for  $t = [n]$ . Let  $S_q$ 

 $=\left\{\left(s_{1}^{(q)}, w_{1}^{(q)}\right), \left(s_{2}^{(q)}, w_{2}^{(q)}\right), \ldots, \left(s_{m}^{(q)}, w_{m}^{(q)}\right)\right\}$ be the signature of a query image. The EMD distance can be converted to an LP optimization problem as follows:

minimize 
$$
\sum_{i=1}^{m_t} \sum_{j=1}^{m_q} f_{i,j} d_{i,j}
$$
  
\n $f_{i,j} \ge 0, 1 \le i \le m_t, 1 \le j \le m_q;$   
\n $\sum_{i=1}^{m_t} f_{i,j} = w_j^{(q)}, 1 \le j \le m_q;$   
\nsubject to  $\sum_{j=1}^{m_q} f_{i,j} = w_i^{(t)}, 1 \le i \le m_t;$   
\n $\sum_{i=1}^{m_t} \sum_{j=1}^{m_q} f_{i,j} = 1.$ 

Similarly, we first need to build a pre-filter table for sub-linear searching time. In the EMD case, we need to filter the images using a quick estimation of EMD without computing it. We utilize an easy-to-compute EMD lower bound as the estimation. The Euclidean distance between the centroids of two signatures with the same total weights is a lower bound on the EMD between them. It is well known that with the same total weights is a<br>
d on the EMD between them. It is<br>  $\begin{array}{ll}\n & \text{where } c \text{ is} \\
d \text{ on the EMD between them. It is} \\
\text{that} \\
\begin{vmatrix}\n m_i & m_q \\
 m_{\text{structure}} & n_{\text{structure}} \\
\text{in } c\n\end{vmatrix} = \sum_{i=1}^{m_i} f_{i,j}^* d_{i,j} \ge \left\| \sum_{i=1}^{m_i} s_i^{(i)} w_i^{(i)} - \sum_{j=1}^{m_q$ 

$$
EMD(S_i, S_q) = \sum_{i=1}^{m_i}\sum_{j=1}^{m_q}f_{i,j}^{\quad \ *} d_{i,j} \geq \left\|\sum_{i=1}^{m_i} s_i^{(t)}w_i^{(t)} - \sum_{j=1}^{m_q} s_j^{(q)}w_j^{(q)}\right\|_2
$$

one-way LSH functions on centroid database one-way LSH functions on centroid database<br>  $\varepsilon$ , then groups the signature ID set on all  $\lambda$  denote the problem (2) by hash tables. After that, the data owner  $(c, U, \tau, V, E)$ . Our design goal is to find a outsources the secure searchable index, encrypted signature database and encrypted secret key image database to the cloud server. image query comes, the data user first generates the signature Q, then compute the generates the signature Q, then compute the<br>centroid  $\xi_q$ . After that, he applies  $\lambda$  LSH optimization functions and one-way function to  $\xi_q$  to  $\eta_{\text{the}}$  the output compute the trapdoor  $T_{\xi_q} = \emptyset \left( K_j, H_j \left( \xi_q \right) \right)$ ,  $j = 1, \dots$ ,  $\lambda$ and sends it ance between the centroids of two<br>
latternes with the same total weights is a<br>
ter bound on the EMD between them. It is<br>  $dD(S_1, S_q) = \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} f_{i,j}^* d_{i,j} \ge \left| \sum_{i=1}^{\infty} s_i^{(i)} w_i^{(i)} - \sum_{j=1}^{\infty} s_j^{(i)} w_j^{(i$ to the cloud server. The cloud server retrieves the corresponding encrypted signatures and<br>sends them to the data user By building the mapping that  $x = y\Lambda - \gamma$ , where  $\Lambda$  is and  $\lambda$ sends them to the data user. By building the secure searchable index, we can achieve the sub-linear filtering complexity for the large scale image database. In the second stage, we will directly compute the underlying distance metric, which is EMD, to compare the similarities of different images. We will explain in the next section. After that, the data owner<br>the secure searchable index, secure and<br>gnature database and encrypted secret key<br>ase to the cloud server. When EMD optithe trapdoor<br>  $T_{\xi_q} = \emptyset (K_j, H_j(\xi_q)), j = 1, ...$ ,  $\lambda$ and sends it want cloud to c<br>
to the cloud server. The cloud server retrieves<br>
the corresponding encrypted signatures and<br>
secure searchable index, we can achieve the sub-li

#### LP Transformation:

The local features always involve in the more complex distance metrics, which need more minimize  $r\mathbf{c}^T \Lambda \mathbf{y} - r\mathbf{c}^T \gamma$ delicate security design. In the second stage, after receiving the encrypted signatures from the cloud in the first stage, the data user decrypts the signatures and computes the distance  $C_{ij}$  between the cluster  $s_i, q_j$ . Then he affect the optima formulates the EMD optimization problem as in Eq.  $(2)$ . The data user wants to leverage the computation power of cloud server to compute computation power of cloud server to compute EMD with privacy protection. Next, we will explain how to make the cloud server securely compare the EMD distance between different images without revealing the sensitive information. For more concise explanation, we use the matrix expression for Eq. (2): will directly compute the underlying distance<br>
metric, which is EMD, to compare the<br>
similarities of different images. We will<br>
Next, we<br>
explain in the next section.<br>
LP Transformation:<br>
The local features always involve

$$
Minimize  $c^T x$   
\nSubject to  $Ux = \tau$  the secret transformation key
$$

 $Vx \le E$  (2)

Similar to the global feature case, we apply<br>one way I SH functions on centroid database<br>vector, where  $1 \le i \le a, 1 \le j \le b$ . We etween the centroids of two<br>
where c is an abx1 distance<br>
and on the EMD between them. It is a where c is an abx1 distance<br>
that<br>
that<br>  $=\sum_{i=1}^{m} \sum_{j=1}^{m} f_{i,j} d_{i,j} \geq \left| \frac{m}{|a|} g_{i,j} w_{ij} \right| = \sum_{j=1}^{m} g_{i,j} w_{ij} \left| \frac{w$ the centroids of two<br>  $Vx \leq E$  (2)<br>
e same total weights is a<br>
EMD between them. It is<br>
abx1 flow vector. U is an 1xab known<br>
structure matrix and V is an  $(x + h) \times ab$ <br>
structure matrix and V is an  $(x + h) \times ab$ <br>
structure matri detained hetewarn the certification is a spheric is an abilitative term in the spheric server s where c is an ab $\times$ 1 distance vector, x is an ab $\times$ 1 flow vector. U is an 1 $\times$ ab known structure matrix and V is an  $(a + b) \times ab$ known structure matrix.  $\tau$  is the minimal total  $Vx \leq E$  (2)<br>where **c** is an ab×1 distance vector, **x** is an<br>ab×1 flow vector. U is an 1×ab known<br>structure matrix and **V** is an  $(a + b) \times ab$ <br>known structure matrix.  $\tau$  is the minimal total<br>weight and  $E = (w_{s_i}, w_{a_j})$  is th  $x \le E$  (2)<br>
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is an 1×ab known<br>
V is an  $(a + b) \times ab$ <br>  $\bar{x}$  is the minimal total<br>  $w_{q_j}$  is the weight<br>  $i \le a, 1 \le j \le b$ . We<br>
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ke c *Vx* ≤ *E* (2)<br>where **c** is an ab×1 distance vector, **x** is an<br>ab×1 flow vector. **U** is an 1×ab known<br>structure matrix and **V** is an  $(a + b) \times ab$ <br>known structure matrix. **r** is the minimal total<br>weight and  $E = (w_{s_i}, w_{q_j})$  i denote the problem (2) by a tuple  $\Psi =$  $Vx \leq E$  (2)<br>where **c** is an ab×1 distance vector, **x** is an<br>ab×1 flow vector. **U** is an 1×ab known<br>structure matrix and **V** is an  $(a + b) \times ab$ <br>known structure matrix. **r** is the minimal total<br>weight and  $E = (w_{s_i}, w_{q_j})$  is t secure and efficient transformation, using a secret key  $K_r$ , to make cloud compute the EMD optimization problem while protecting the input and output privacy. It means t cloud solves the randomly transformed EMD<br>optimization problem without knowing the optimization problem without knowing the input distance vector  $c$ , weight vector  $E$  and the output flow vector x. After that, we also want cloud to compute the EMD distance and want cloud to compute the EMD distance and<br>sort the results. Our secure transformation contains two main steps. First of all, to protect the output privacy, we perform the affine structure matrix and **V** is an  $(a + b) \times ab$ <br>known structure matrix.  $\tau$  is the minimal total<br>weight and  $E = (w_{s_i}, w_{q_j})$  is the weight<br>vector, where  $1 \le i \le a, 1 \le j \le b$ . We<br>denote the problem (2) by a tuple  $\Psi =$ <br> $(c, U, \tau, V, E)$ ab non-singular matrix and  $\gamma$  is an ab  $\times$  1 vector. The original problem is transformed to: vector. The vector, where  $1 \le i \le a, 1 \le j \le b$ . We<br>
denote the problem (2) by a tuple  $\Psi =$ <br>  $(c, U, \tau, V, E)$ . Our design goal is to find a<br>
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efficient transformation, using a<br>
efficient transformation, using a<br>  $K_r$ , to make cloud compute the<br>
nization problem while protecting<br>
d output privacy. (2)<br>ctor, **x** is an<br> $1 \times ab$  known<br> $(a + b) \times ab$ <br>minimal total<br>the weight<br> $\exists j \leq b$ . We<br>u tuple  $\Psi =$ <br>is to find a weight and  $E = (w_{s_i}, w_{q_i})$  is the weight 1  $(2)$ <br>
1  $\rightarrow$  1  $\times$  2  $\rightarrow$  1  $\times$  2  $\rightarrow$  2  $\rightarrow$  2  $\rightarrow$  2  $\rightarrow$  2  $\rightarrow$  3  $\rightarrow$  4  $\rightarrow$  4  $\rightarrow$  4  $\rightarrow$  4  $\rightarrow$  4  $\rightarrow$  9  $\rightarrow$  4  $\rightarrow$  9  $\rightarrow$  9  $\rightarrow$  1  $\rightarrow$  1 secure and efficient transformation, using a<br>secret key  $K_r$ , to make cloud compute the<br>EMD optimization problem while protecting<br>the input and output privacy. It means that the  $(a + b) \times ab$ <br>inimal total<br>the weight<br> $\leq b$ . We<br>tuple  $\Psi =$ <br>s to find a<br>onn, using a<br>ompute the<br>protecting<br>ans that the<br>protecting<br>ans that the<br>rector **E** and<br>at, we also<br>listance and<br>ant, we also<br>listance and<br>nsformation<br>1, tructure matrix.  $\tau$  is the minimal total<br>and  $E = (w_{s_i}, w_{q_j})$  is the weight<br>where  $1 \le i \le a, 1 \le j \le b$ . We<br>the problem (2) by a tuple  $\Psi =$ <br> $\gamma, E$ ). Our design goal is to find a<br>and efficient transformation, using a<br>gy  $K_r$  $\leq j \leq b$ . We<br>
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minimize 
$$
\mathbf{c}^T \mathbf{\Lambda} \mathbf{y} - \mathbf{c}^T \mathbf{\gamma}
$$
  
subject to  $\mathbf{U} \mathbf{\Lambda} \mathbf{y} = \tau + \mathbf{U} \mathbf{\gamma}$ ,  
 $\mathbf{V} \mathbf{\Lambda} \mathbf{y} \le \mathbf{E} + \mathbf{V} \mathbf{\gamma}$  (3)

Next, we multiply the  $((a + b) \times (a + b))$ generalizes permutation G to the inequality constraints to protect E and multiply a real positive value r to protect optimal value. We transform the problem (3) to another problem:

subject to 
$$
U\Lambda y = \tau + U\gamma
$$
,  
\n $GV\Lambda y \leq G(E + V\gamma)$  (4)

,  $q_i$ . Then he affect the optimal solution, the final Because the constant term  $rc^T\gamma$  does not transformation problem can be formed as:

minimize 
$$
\mathbf{c}^{\prime T} \mathbf{y}
$$
  
bject to  $\mathbf{U}'\mathbf{y} = \tau'$ ,  
 $\mathbf{V}'\mathbf{y} \leq \mathbf{E}'$ 

Minimize  $c^T x$  the secret transformation key Where  $c^T = rc^T \Lambda$ ,  $U' = U \Lambda$ ,  $\tau' = \tau +$ vector. The original probe<br>
minimize  $c^T A y - c^T \gamma$ <br>
subject to  $U A y = \tau + U \gamma$ ,<br>  $V A y \le E + V \gamma$  (3)<br>
Next, we multiply the  $((a + b) \times (a + b))$ <br>
constraints permutation G to the inequality<br>
constraints to protect E and multiply a re problem has similar structure to the problem has similar structure to the original problem (2). We use $\varphi_{K_r} = (c', 0', \tau', 0', E')$ to denote the secure transformed problem, where denote data soluted to:<br>  $-cT\gamma$ <br>  $+ U\gamma$ , (3)<br>
the inequality<br>
multiply a real<br>
imal value. We<br>
other problem:<br>
(4)<br>  $\gamma$  does not<br>
final<br>
formed as:<br>
(5)<br>  $\gamma' = \tau +$ <br>
(5). This<br>
b the original<br>  $U', \tau', V', E'$ ) to<br>
problem, where  $U\gamma$ ,  $V' = GV\Lambda$ ,  $E' = G(E + V\gamma)$ . This + b) × (a + b))<br>
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multiply a real<br>
timal value. We<br>
mother problem:<br>
(4)<br>  $\gamma$  does not<br>
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formed as:<br>
(5)<br>  $\tau' = \tau +$ <br>  $\gamma$ ). This<br>
o the original<br>  $J'$ ,  $\tau'$ ,  $V'$ ,  $E'$ )to<br>
problem, where (v)<br>
(v)  $\times$  (a + b))<br>
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(5)<br>  $\tau$  +<br>
This<br>
e original<br>  $\tau$ ,  $V$ ,  $E$ ) to<br>
blem, where  $\times$  (a + b))<br>inequality<br>ply a real<br>value. We<br>problem:<br>(4)<br>s not<br>ed as:<br>(5)<br>+<br>is<br>original<br>, V', E')to<br>em, where  $(a + b)$ <br>
equality<br>  $(a + b)$ <br>  $(c + b)$ <br>  $(c + b)$ <br>  $(c + b)$ <br>  $(c + b)$ <br>  $(d + b)$ <br>
oblem:<br>
(4)<br>
ot<br>  $s$ <br>  $\frac{1}{2}$ <br>  $K_r = (G, \Lambda, \gamma, r)$ . The whole transformation process is illustrated in Algorithm 1. After solving the EMD optimization problem, we also want cloud to compute EMD as in Eq. (1) and sort the results. Observe that the numerator of EMD equation is the optimal values of the original problem (2) and the denominator is the sum of the elements of optimal solution. However, the cloud server solves the transformed problem (5) instead of the original one for security reason. The difference of optimal value between original problem and transformed problem is a constant term $rc^{T}\gamma$ , which can be computed by data users before outsourcing. We also want the sum of elements of optimal solutions between original problem transformed problem differs in a constant term after the affine mapping  $x = \Lambda y - \gamma$ . So, we apply one additional constraint when constructing  $\Lambda$  such that for  $\Lambda^{-1}$  each column's sum is one. Then we can derive that  $\sum_{i=1}^{a} \sum_{j=1}^{b} x_{ij} = \sum_{i=1}^{a} \sum_{j=1}^{b} y_{ij}$  –  $\sum_{k=1}^{ab} \gamma_k$ . The data user can compute  $\sum_{k=1}^{ab} \gamma_k$  before outsourcing. When outsourcing the secure transformed problems to cloud server, the data users also send two offset constants  $rc^T \gamma$  and $\sum_{k=1}^{ab} \gamma_k$ to cloud server. Then the cloud server can solve the transformed problems and use these constants to compute the right EMD in an order preserving way. The offset constants themselves will reveal little information. After that, the cloud server sorts the results and returns the more accurate ranked order encrypted images.

## Global Feature vs. Local Feature

Under our two stages similarity image search system, the global and local feature-based solutions have certain differences. The global feature-based solution only needs one round communication between data user and cloud server to retrieve the ranked order encrypted images, while local feature based solution needs two rounds. There are two reasons. The first reason is that the underlying distance metric for local feature based solution, which is EMD, has more complicated problem structure and needs more delicate security design. Our encryption method should preserve the LP problem's structure and make sure the cloud server can solve the right optimal solution. We need to compute the distance vector $c_{zq}^{T}$  for each retrieved and querying signature pair first, then encrypt this distance vector using transformation key  $\Lambda$ , r. We cannot outsource these two operations to cloud server at the same time without revealing the distance vector and the transformation key. The second reason is that there is a trade-off between communication round and security. For one round solution, every feature vector encrypted use the same invertible matrix W. It is vulnerable to know plaintext attack, when cloud server has certain number of pairs of plaintexts and ciphertext feature vectors, it can derive the invertible matrix W. For two rounds solution, the data owner adopts the standard encryption techniques, which can resist the known plaintext attack. So, we favor the two rounds design for our system, which also has higher scalability for different traditional and underlying distance metrics.

User:

In this application user is a module, here user should register with the application then login. After user successful login he can perform some operations, such as generate Key, search Images, request Status and logout

#### Owner:

In this application user is a module, here user should register with the application then login. After user successful login he can perform some operations, such as generate Key, upload Image, view Request and logout.

### Cloud:

Here cloud can directly login with the application and after cloud successful login he can perform some operations such as view all images and view all users requests and logout.

## OUTPUTS :







#### CONCLUSION :

In this paper, we investigate similarity search  $\frac{S_y}{R}$ for encrypted images in secure cloud<br>computing. Concretely, we intro-duce a Thermodyn computing. Concretely, we intro-duce a clustering improvement method and give the design method of the hierarchical index tree.  $\frac{1}{3}$ With these two techniques, SEI can efficiently With these two techniques, SEI can efficiently<br>
perform the retrieval process and achieve high<br>
Isoured of Imperi accuracy based on features extracted by the CNN model. Further, we consider un trusted image users in SEI and hence propose a similarity calculation method with limited 4. key-leakage. We also give strict security analysis and conduct experiments on world dataset, which indicate that SEI is secure and feasible in practice. racy based on features extracted by the<br>V model. Further, we consider un trusted<br>ge users in SEI and hence propose a<br>larity calculation method with limited<br>leakage. We also give strict security<br>ysis and conduct experiments

#### FUTURE ENHANCEMENT :

Efficiency Optimization: Explore techniques like selective encryption or partial feature indexing to further reduce computational overhead during the search process, especially for very large datasets.

Advanced Searchable Encryption: Investigate the integration of emerging SE schemes that offer additional functionalities like multi-keyword search or fuzzy search, allowing for more complex user queries. iciency Optimization: Explore techniques<br>
exing to further reduce computational<br>
rhead during the search process, especially<br>
very large datasets.<br> **vanced** Searchable Encryption:<br>
estigate the integration of emerging SE<br>

Privacy Amplification: Privacy Amplification: Implement techniques like random masking homomorphic encryption to further strengthen privacy guarantees against potential key leakage scenarios. ike random masking or<br>encryption to further strengthen<br>antees against potential key<br>ios.<br>**ata Management:** Develop

Dynamic Data Management: Develop mechanisms for efficient insertion, deletion, and update of image data while maintaining the integrity and searchability of the encrypted index structure.

Scalability Evaluation: Conduct extensive scalability tests on even larger and more diverse image datasets to assess the system's performance under high load. For deletion, deletion,<br>
update of image data while maintaining<br>
integrity and searchability of the encrypted<br>
Ex structure.<br> **lability Evaluation:** Conduct extensive<br>
ability tests on even larger and more<br>
rese image data

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